

The Method of Solv Cands

Ben De Bonet + Ohad Karon

Φ satisfying prop formula in CNF

A solution is $\mu: \nu \rightarrow \{0, 1\}$ s.t. $\mu \models \Phi$.

Can we add a solution by forcing?

$H(\Phi)$ control forcing for this

| | | | |
|----|--------------------|--------------------|-----|
| I | $c_0 \models \Phi$ | $c_1 \models \Phi$ | ... |
| II | $d_0 \in c_0$ | $d_1 \in c_1$ | ... |

or moves

II must not push $d_i, \neg d_i$ for any i .

$G(\Phi, p)$, p = partial assignment.

Cloud game for player II, here determined.

$p \in H(\Phi)$ if $\exists \Gamma \vdash G(\Phi, p)$.

$a \leq p$ if $q \geq p$

$H(\Phi)$ may collapse controls etc

Can we add an assignment w/o causing chaos, or by a point $P \in K$, K = new class of forcing e.g. few proper

$$P_0 = H(\Phi)$$

$P_1 = \text{Fix } \theta, M \prec H_\theta \text{ c.t.b.l., } p \in H(\Phi)$

| | |
|---|------------------|
| I | $c \models \Phi$ |
| | $d \in c$ |

$$P_{n+1} = P_n \cup \{\# \}$$

$$\begin{cases} P_0 = p \\ V \\ P_1 \\ V \\ P_2 \\ \vdots \\ P_{n-1} \end{cases}$$

Player 1 can play a den $D \in P_0$, $D \subseteq M$

(2)

Player 2 picks $a \in D$, $p_n \leq p_{n-1}$, a

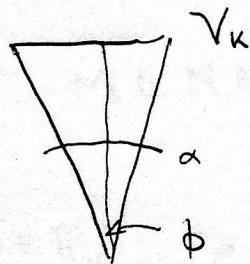
$$\text{Hull}(M, \{a\}) \cap w_i = M \cap w_i.$$

M is a good model if II has a very strategy to L
 $G(\phi, M, p) \quad \forall p \in M \cap P_0.$

$\cdot p \in P_1: p = (M_p, w_p) \quad |M_p| \leq 1.$

$\bullet M$ good n.h.t., esp II $\uparrow G(\phi, M, \omega_p).$

we want this to "converge"



$$\varepsilon = \{\alpha: V_\alpha \supseteq V_n\}$$

$$V_\alpha \supseteq V_n$$

$$\alpha \in \varepsilon, \alpha \text{-models}$$

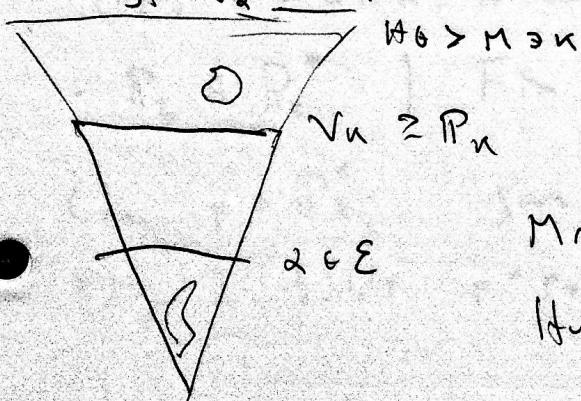
M is an α -model if

1. $M \subset \hat{M} \models ZFC^-$ for \hat{M} = transitive closure
 not ctbl

2. $\hat{M} = \text{Hull}(M, V_\alpha) =$

$$\{f(\bar{x}): f \in M, \bar{x} \in V_\alpha\}$$

3. $V_\alpha \in M$.



$$M \cap n \in \alpha.$$

$$\text{Hull}(M, V_\alpha)$$

(5)

$$HOM(M, V_\alpha) \cap N = \emptyset$$

$$\pi_{\text{loc}}$$

$$\overline{HOM(M, V_\alpha)}$$

$$\pi[M] \in \text{an } \alpha\text{-model}$$

"

$$M_{T_\alpha}$$

$$p \in M$$

$$\overset{\text{"}}{(M_p, w_p)}$$

$$C_\alpha = \{M : M \text{ a model}\}$$

$$C = \bigcup C_\alpha$$

Preconditions

$$p \in P_n^* \text{ if } p = (M_p, w_p)$$

$$- w_p \in H(\phi)$$

- M_p is a finite ϵ -chain of models in C .

$$M \in N \in M_p \rightarrow \lambda_M \not< \lambda_N.$$

"the κ of the model"

$$p \in P_n^*, \quad p \upharpoonright \lambda = (\{M \in M_p : \lambda_M < \lambda\}, w_p)$$

$$q \leq p \text{ iff } w_q \geq w_p \text{ & } \forall M \in M_p \exists N \in M_q \quad \lambda_M = \lambda_N$$

and $M \leq_{w_1} N$. (cannot reverse w_1)

$$S_M = M \cap w_1.$$

We define P_λ , $\lambda \in \epsilon$ by ~~induction~~.

- $(P_\lambda, \zeta \in \epsilon \cap \lambda)$ is defined

- $P_\lambda \in P_\zeta^* \upharpoonright \text{Fix } \lambda \in \epsilon$

Given $p \in P_\lambda^*$, $\beta \in G_\lambda(p)$

Player II builds $p = p_0 \succ p_1 \succ p_2 \dots$ legal w

(4)

| | |
|-----------|---------------------------------|
| <u>I</u> | <u>Q_n</u> ✓ question |
| <u>II</u> | |

I Q_n: C_n ∈ \$

II l_n ∈ C_n p_n = (M_{p_n}, w_{p_n}, {l_n})

I Q_n: (M, D), M ⊂ M_{p_n}, D ⊂ M den ⊂ P_{D_n}

II prakt q ∈ D, p_{n+1} ∈ p_n, q
P_{D_n}*

$$\text{Hull}(M, \{q\}) \cap \omega_1 = M \cap \omega_1$$

$$(\delta_{\text{Hull}(M, \{q\})} = \delta_M)$$

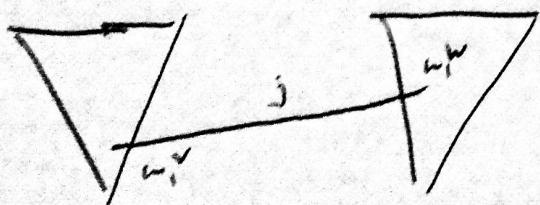
p ∈ P_{D_n} iff II aus G_{D_n}(p).

(Kern) AS-constant.

φ ist AS-constant & für every int p_n

J_n √^{col(w_n, u_n)} H_μ √^{col(w_n, u_n)} μF + ∃ j: V → W

$$\text{car } j = w_i^V \circ j(S)$$



∃ A F_j(φ) s.t. j[n] ⊂ μ.